



## Dominators for Multiple-objective Quasiconvex Maximization Problems

EMILIO CARRIZOSA<sup>1,\*</sup> and FRANK PLASTRIA<sup>2</sup>

<sup>1</sup>Facultad de Matemáticas, Universidad de Sevilla, C/Tarfia s/n, 41012 Sevilla, Spain

E-mail: [ecarriz@cica.es](mailto:ecarriz@cica.es)

<sup>2</sup>Department of Management Informatics, Vrije Universiteit Brussel, Pleinlaan, 2, B-1050 Brussels, Belgium (E-mail: [Frank.Plastria@vub.ac.be](mailto:Frank.Plastria@vub.ac.be))

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**Abstract.** In this paper we address the problem of finding a dominator for a multiple-objective maximization problem with quasiconvex functions. The one-dimensional case is discussed in some detail, showing how a Branch-and-Bound procedure leads to a dominator with certain minimality properties. Then, the well-known result stating that the set of vertices of a polytope  $S$  contains an optimal solution for single-objective quasiconvex maximization problems is extended to multiple-objective problems, showing that, under upper-semicontinuity assumptions, the set of  $(k - 1)$ -dimensional faces is a dominator for  $k$ -objective problems. In particular, for biobjective quasiconvex problems on a polytope  $S$ , the edges of  $S$  constitute a dominator, from which a dominator with minimality properties can be extracted by Branch-and Bound methods.

**Key words:** Multiple-objective problems; Quasiconvex maximization; Dominators

### 1. Introduction

Given a nonempty closed subset  $S$  of  $\mathbb{R}^n$  and a function  $F : S \subset \mathbb{R}^n \rightarrow \mathbb{R}^k$ , define the *multiple-objective* problem  $(P[F; S])$ ,

$$\max_{x \in S} F(x), \quad (P[F; S])$$

which seeks those alternatives maximizing simultaneously the components  $F_1, F_2, \dots, F_k$  of  $F$ , [7, 28, 31].

Although the term simultaneous maximization is not uniquely defined, it customarily means finding the set  $\mathcal{E}[F; S]$  of *efficient* or *Pareto-optimal* solutions to  $(P[F; S])$ ,

$$\mathcal{E}[F; S] = \{x \in S : \text{no } y \in S \text{ verifies } F_i(y) \geq F_i(x) \forall i = 1, 2, \dots, k \\ \text{with at least one inequality strict}\}$$

In general  $\mathcal{E}[F; S]$  lacks many desirable properties such as being connected or closed, and this seems to be quite often the case and not only in pathological

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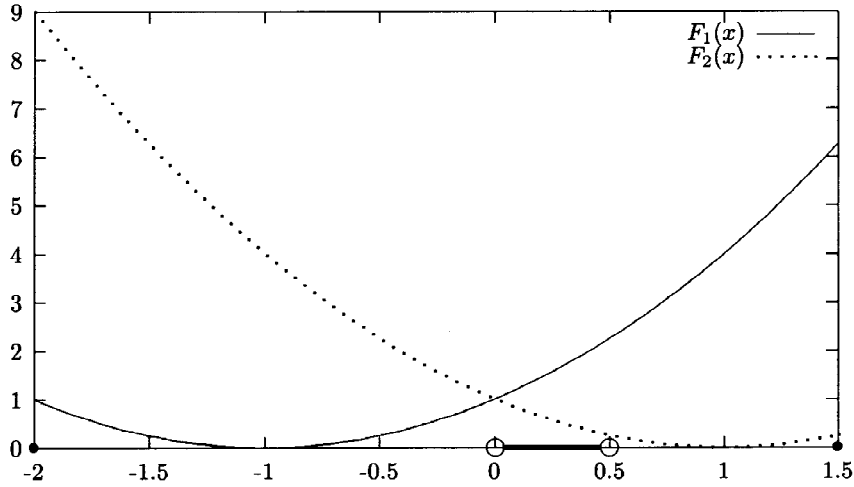


Figure 1. Biobjective convex maximization.

examples: take, for instance, the biobjective convex maximization problem in one variable ( $n = 1$ ,  $k = 2$ ) with  $F(x) = ((x + 1)^2, (x - 1)^2)$  and  $S = [-2, 1.5]$ , plotted in Figure 1.

Since  $F(-2) \geq F(x) \forall x \in ]-2, 0]$ , with at least one inequality strict, and  $F(1.5) \geq F(x) \forall x \in [0.5, 1.5[$ , with at least one inequality strict too, it follows that the set of Pareto-optimal points must be contained in  $\{-2\} \cup ]0, 0.5[ \cup \{1.5\}$ . In fact, it is readily seen from the plot that

$$\mathcal{E}[F; S] = \{-2\} \cup ]0, 0.5[ \cup \{1.5\},$$

which is a disconnected non-closed set. See following sections and also e.g. [3] for other instances.

Moreover, although there exist procedures to check whether a given point is efficient or not, e.g. [7, 31], an algorithm to construct  $\mathcal{E}[F; S]$  is only available for a few classes of problems, such as multiple-objective linear problems, [28].

This drawback has been overcome in the literature by means of two strategies: either  $\mathcal{E}[F; S]$  is sought, but, due to the inability for obtaining it, an approximation (sometimes with unknown degree of precision) is provided, e.g. [8, 18], or else the concept of efficiency is relaxed and replaced by a manageable surrogate of it.

In this paper we follow the second approach by using the concept of *dominator*, [5, 16, 21, 30], also called *weak kernel*, e.g. in [31] which is defined as any subset  $S_0 \subset S$  such that, for any feasible  $x \notin S_0$ ,  $S_0$  contains a feasible alternative at least as good as  $x$  with respect to all objectives. See Section 2 for a formal definition.

It should be remarked that this concept is not only useful as a surrogate of the idea of Pareto-efficiency, but also as a tool in the resolution of some single-objective problems. Indeed, some of the most popular optimization methods for single-objective problems of the form

$$\max_{x \in S} \Psi(x) \tag{1.1}$$

require the feasible region  $S$  to be bounded. Such is the case, among others, of the Branch and Bound methods for global optimization, e.g. [15], which, in their simplest version, require, as pre-processing, the construction of a bounded polyhedron  $P$  (usually a hyper-rectangle, or a simplex) including either the whole feasible region, or, at least a bounded subset  $S_0 \subset S$  known to contain an optimal solution. Moreover, the speed of convergence of the procedure is known to deteriorate with the volume of  $P$ , so  $P$  should be as small as possible in order to obtain reasonable computation times.

How to construct  $P$  will depend, of course, on the specific properties of the problem at hand. In particular, if (1.1) has the form

$$\max_{x \in S} \Phi(F(x)), \tag{1.2}$$

for some  $\Phi : F(S) \rightarrow \mathbb{R}$  componentwise non-decreasing, then it is well known that, if (1.1) has optimal solutions, then any dominator for the multiple-objective problem  $\max_{x \in S} F(x)$  also contains optimal solutions for (1.1), [21]. In other words, we can take as  $S_0$  any bounded dominator for the multiple-objective problem, and as  $P$  any superset of  $S_0$  with the required geometry.

This property has been successfully exploited, among others, in [5, 21, 22, 30] for problems of Linear Regression and Continuous Location, in which the globalizing function  $\Phi$  is an arbitrary non-decreasing function and the function  $F$  is componentwise concave. Our aim here is to address the (harder) problem in which the function  $F$  is componentwise (quasi)-convex, showing as main result (Proposition 19) that, under upper-semicontinuity assumptions, the search of a dominator can be restricted to the  $(k - 1)$ -dimensional faces of  $S$ .

The rest of this paper is structured as follows. In Section 2 we formally introduce the concept of dominators and discuss some general properties. These properties are used in Section 3 to address the one-dimensional case, for which dominators with certain minimality properties can be obtained.

Section 4 is devoted to show that, for multiple-objective multi-dimensional problems, one can construct dominators contained in low dimensional faces of the polytope  $S$ .

The paper ends with an application of these results to the construction of a dominator for a biobjective problem in Continuous Location. The reader is referred also to [25] for another successful application of the technique developed in this paper.

## 2. Dominators

Defining for each  $x \in S$  the upper level set at  $x$  of  $F$  on  $S$ ,  $\mathcal{S}^{\geq}(x)$  as

$$\mathcal{S}^{\geq}(x) = \{y \in S : F_i(y) \geq F_i(x) \text{ for all } i = 1, 2, \dots, k\},$$

the set  $\mathcal{E}[F; S]$  of efficient solutions may be defined by

$$\begin{aligned}\mathcal{E}[F; S] &= \{x \in S : \text{If } y \in \mathcal{S}^{\geq}(x) \text{ then } x \in \mathcal{S}^{\geq}(y)\} \\ &= \{x \in S : \text{If } y \in \mathcal{S}^{\geq}(x) \text{ then } F(x) = F(y)\}\end{aligned}$$

DEFINITION 1. A set  $S^* \subset S$  is said to be a dominator for  $(P[F; S])$  iff for each  $x \in S$  there exists some  $x^* \in S^*$  which has, componentwise, a value not smaller than  $x$ . In other words,  $S^*$  is a dominator iff

$$(\forall x \in S) \exists x^* \in \mathcal{S}^{\geq}(x) \cap S^*$$

Hereafter, the class of dominators for  $(P[F; S])$  will be denoted by  $\mathcal{D}[F; S]$ .

A direct consequence of the definition is the following:

PROPOSITION 2. One has

1.  $S \in \mathcal{D}[F; S]$ . In particular,  $\mathcal{D}[F; S]$  is nonempty.
2. If  $D \in \mathcal{D}[F; S]$  and  $D^* \subset D \subset D^* \subset S$ , then  $D^* \in \mathcal{D}[F; S]$ .
3. For any class  $\{S_j : j \in J\}$  of nonempty sets in  $\mathbb{R}^n$ ,

$$\text{If } S_j^* \in \mathcal{D}[F; S_j] (\forall j \in J) \text{ then } \bigcup_{j \in J} S_j^* \in \mathcal{D}\left[F; \bigcup_{j \in J} S_j\right]$$

4. For any class  $\{S_j : j \in J\}$  of nonempty sets in  $\mathbb{R}^n$ ,

$$\bigcap_{j \in J} \mathcal{D}[F; S_j] \subset \mathcal{D}\left[F; \bigcup_{j \in J} S_j\right]$$

5. If  $D \in \mathcal{D}[F; S]$ , then

$$\mathcal{D}[F; D] \subset \mathcal{D}[F; S].$$

By Proposition 2, the class  $\mathcal{D}[F; S]$  is nonempty since the whole feasible set  $S$  is one of its elements. However  $S$  does not seem to be the most appropriate dominator since it possibly contains (too) many dominated alternatives, being too far from the ideal aim of a smallest possible dominator.

PROPOSITION 3. Suppose each  $F_j$  is upper-semicontinuous on  $S$ , then any class of compact nested dominators is closed under intersections. In other words: if  $(I, \leq)$  is a totally ordered set, and  $\{D_i\}_{i \in I}$  is a class of compact dominators with  $D_i \subset D_j$ ,  $j \in I$ ,  $i \leq j$ , then

$$\bigcap_{i \in I} D_i \in \mathcal{D}[F; S].$$

*Proof.* Take any  $x \in S$ . By the upper-semicontinuity of the functions  $F_j$ , all upper level sets  $\{y \in S : F_j(y) \geq F_j(x)\}$  are closed, so their intersection  $\mathcal{S}^{\geq}(x)$  is also closed. By the definition of dominators and their compactness, it follows for each  $i \in I$  that  $\mathcal{S}^{\geq}(x) \cap D_i$  is a nonempty compact set, thus  $\{\mathcal{S}^{\geq}(x) \cap D_i\}_{i \in I}$  constitutes a

class of nested compact sets. By compactness their intersection (i.e.,  $\mathcal{S}^\geq(x) \cap \bigcap_{i \in I} D_i$ ) is nonempty.

However, it is evident that the whole class  $\mathcal{D}[F; S]$  is not closed under intersections (take constant functions  $F_1, \dots, F_k$ , then any singletons  $\{x\}, \{y\} \subset S$  are dominators, with empty intersection). Hence, a unique smallest dominator is unlikely to exist. We then relax the idea of smallest dominator by introducing the concept of (weak) minimal dominators. First define for each  $x \in S$  the strict upper level set of  $F$  on  $S$ ,  $\mathcal{S}^>(x)$  as

$$\mathcal{S}^>(x) = \{y \in S : F_i(y) > F_i(x) \text{ for all } i = 1, 2, \dots, k\}.$$

DEFINITION 4. A dominator  $S^*$  is said to be minimal for  $(P[F; S])$  iff no proper subset of  $S^*$  belongs to  $\mathcal{D}[F; S]$ . In other words,  $S^* \subset S$  is minimal iff

$$(x, y \in S^*, x \neq y) \Rightarrow x \notin \mathcal{S}^\geq(y)$$

A dominator  $S^* \subset S$  is said to be weak minimal for  $(P[F; S])$  iff

$$(x, y \in S^*) \Rightarrow x \notin \mathcal{S}^>(y)$$

The class of minimal (respectively weak minimal) dominators for problem  $(P[F; S])$  will be denoted by  $\mathcal{D}_M[F; S]$  (respectively  $\mathcal{D}_{WM}[F; S]$ ).

As a simple illustration of the concepts, consider the 2-dimensional 2-objective optimization problem  $\max_{x \in S} F(x)$ , depicted in Figure 2, where the feasible region  $S$  is the polyhedron in  $\mathbb{R}^2$  with vertices  $a = (0, -3)$ ,  $b = (4, -1)$ ,  $c = (4, 0)$ ,  $d = (0, 3)$ , and  $F$  is given by

$$F_1(x_1, x_2) = x_1$$

$$F_2(x_1, x_2) = |x_2|$$

Then, the Pareto optimal set is given by

$$\mathcal{E}[F; S] = \{d\} \cup [a, b],$$

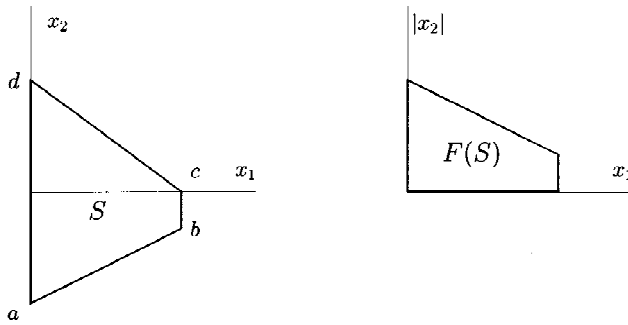


Figure 2.  $S$  and  $F(S)$ .

only two minimal dominators exist, namely

$$\begin{aligned} S_1 &= [a, b] \\ S_2 &= ]a, b] \cup \{d\}, \end{aligned}$$

whereas the polygonal  $S_3$ ,

$$S_3 = \{d\} \cup [a, b] \cup [b, c]$$

is also weak minimal.

We observe in this example that the two minimal dominators are proper subsets of  $\mathcal{E}[F; S]$ . This result is more general, as stated in the following:

**PROPOSITION 5.** *Suppose that  $S$  is compact and each  $F_i$  is upper semicontinuous on  $S$ . Then*

1.  $\mathcal{E}[F; S]$  is a weak minimal dominator.
2. Minimal dominators exist.
3.  $\mathcal{E}[F; S] = \bigcup_{S^* \in \mathcal{D}_M[F; S]} S^*$ .

*Proof.* By the upper-semicontinuity assumption, for each  $x \in S$  the set  $\mathcal{S}^\geq(x)$  is compact. Hence, by Theorem 6 of Chapter 2 of [31]  $\mathcal{E}[F; S]$  is a dominator, which, by construction, is also weak minimal. Hence 1 holds.

To show 2, define on  $\mathcal{E}[F; S]$  the equivalence relation

$$\rho = \{(x, y) \in \mathcal{E}[F; S] \times \mathcal{E}[F; S] : F(x) = F(y)\}.$$

Taking exactly one element in every equivalence class, we obtain a set  $S^*$  which is, by construction, a minimal dominator. Indeed, it is a dominator because  $\mathcal{E}[F; S]$  is a dominator, as shown in Part 1. Moreover it is minimal: if there exists some dominator  $M \subset S^*$ ,  $M \neq S^*$ , for any  $x \in S^* \setminus M$  there would exist some  $y \in M$  with  $F(y) \geq F(x)$ . But by construction of  $S^*$  we would have  $F(y) \neq F(x)$  contradicting the fact that  $x$  is efficient. Hence, minimal dominators exist.

For Part 3, we first show that every efficient point is in some minimal dominator: let  $x^* \in \mathcal{E}[F; S]$ , and construct a subset  $S^*$  of  $\mathcal{E}[F; S]$  taking exactly one element of every equivalence class (with respect to the equivalence relation  $\rho$  above),  $x^*$  being the element chosen from its equivalence class. Using the reasoning above, it is seen that  $S^*$  is a minimal dominator, and  $x^* \in S^*$ .

Finally to show that any minimal dominator is included in the efficient set, take  $x^* \in S^*$ , for some  $S^* \in \mathcal{D}_M[F; S]$ , and assume  $x^* \notin \mathcal{E}[F; S]$ . Then, there exists some  $y \in S$  with  $F(y) \geq F(x)$ , and at least one inequality strict. Since  $S^* \in \mathcal{D}_M[F; S]$ , there must exist some  $y^* \in S^*$  with  $F(y^*) \geq F(y) \geq F(x^*)$ , thus the set  $S^* \setminus \{x^*\}$  will also be a dominator, contradicting the minimality of  $S^*$ . Hence,  $x^* \in \mathcal{E}[F; S]$ .  $\square$

**REMARK 6.** The upper-semicontinuity assumption is needed in order to guarantee the nonvoidness of  $\mathcal{D}_{WM}[F; S]$ , as the following counterexample shows: Let  $S \subset \mathbb{R}^2$  be the triangle whose endpoints are  $(-1, 0)$ ,  $(1, 0)$ ,  $(0, 1)$ , and let  $F_1 : S \rightarrow \mathbb{R}$  be

defined as  $1/(1-x_2)$  on the relative interior of the two top-edges, and zero elsewhere. Since

$$\lim_{\substack{(x_1, x_2) \rightarrow (0, 1), \\ (x_1, x_2) \in bd(S)}} F_1(x_1, x_2) = +\infty,$$

the maximum of  $F_1$  on  $S$  is not attained, thus any  $D \in \mathcal{D}[F_1; S]$  must contain a sequence of boundary points converging to  $(0, 1)$ , implying that  $D$  contains points  $x, y$  with  $F_1(x) > F_1(y)$ . Hence, no weak minimal dominator exists.  $\square$

### 3. Multiple-objective one-dimensional problems

In this section we address the multiple-objective problem  $(P[F; S])$  when  $S$  is given as a finite union of compact intervals in  $\mathbb{R}$ , and each  $F_i$  is quasiconvex on each interval. We first discuss some properties of one-dimensional single-objective quasiconvex minimization problems, which are then used to tackle  $(P[F; S])$ , first when  $S$  reduces to a single compact interval and then in the general case. For the basic properties of quasiconvex functions we refer the reader to [1].

#### 3.1. SINGLE-OBJECTIVE QUASICONVEX MINIMIZATION PROBLEMS ON AN INTERVAL

Let  $I \subset \mathbb{R}$  be a nonempty compact interval, and let  $g : I \rightarrow \mathbb{R}$  be quasiconvex. We will denote by  $\text{cl}_I g$  the closure of  $g$  relative to  $I$ , namely

$$\begin{aligned} \text{cl}_I g(x) &= \inf \{ t : \exists \{x_r\}_r \subset I, \text{ such that } x_r \rightarrow x, g(x_r) \rightarrow t \} \\ &= \liminf_{x_r \rightarrow x} g(x_r) \end{aligned} \quad (3.3)$$

LEMMA 7. *One has:*

1.  $g(x) \geq \text{cl}_I g(x)$  for all  $x \in I$ .
2.  $\inf_{x \in I} g(x) = \inf_{x \in I} \text{cl}_I g(x)$ .
3.  $\text{cl}_I g$  is quasiconvex and lower-semicontinuous.
4. The set  $\arg \min_{x \in I} \text{cl}_I g(x)$  of optimal solutions to  $\min_{x \in I} \text{cl}_I g(x)$  is a nonempty compact subinterval of  $I$ .

*Proof.* 1 to 3 immediately follow from the definition of quasiconvexity and (3.3).

By the lower semicontinuity of  $\text{cl}_I g$ , the set  $\arg \min_{x \in I} g(x)$  is compact and nonempty; since  $\text{cl}_I g$  is also quasiconvex, it follows that  $\arg \min_{x \in I} \text{cl}_I g(x)$  is also convex, thus it is a compact interval, and Part 4 follows.  $\square$

We recall that a function  $g$  is said to be *semistrictly quasiconvex*, [1], if it satisfies the following:

$$\left. \begin{array}{l} g(a) < g(b) \\ c \in ]a, b[ \end{array} \right\} \Rightarrow g(c) < g(b)$$

The next lemma shows that, due to the quasiconvexity of  $g$ , the behavior of  $g$  and

$\text{cl}_1 g$  are closely related, the relationship being stronger for semistrictly quasiconvex  $g$ :

LEMMA 8. *Let  $x^* \in \arg \min_{x \in I} \text{cl}_1 g(x)$ , and let  $z_1, z_2 \in I$  such that  $z_1 \in ]x^*, z_2[$ . One has:*

1.  $g(z_1) \leq g(z_2)$ .
2. *If  $g$  is also semistrictly quasiconvex and  $g(z_1) = g(z_2)$ , then  $]x^*, z_2[ \subset \arg \min_{x \in I} g(x)$ .*

*Proof.* By definition of  $\text{cl}_1 g$  and Part 2 of Lemma 7, one can take a sequence  $\{x_r\}$  in  $I$  converging to  $x^*$  such that  $\inf_r g(x_r) = \inf_{x \in I} g(x) = \text{cl}_1 g(x^*)$ .

Since  $z_1 > x^*$ , there exists  $r_0$  such that  $x_r < z_1$  for all  $r \geq r_0$ , thus

$$z_1 \in ]x_r, z_2[ \text{ for all } r \geq r_0$$

Given  $r \geq r_0$ , if it were the case that  $g(z_2) < g(z_1)$ , then

$$\begin{aligned} g(z_2) &< g(z_1) \\ &\leq \max\{g(z_2), g(x_r)\} \end{aligned}$$

Hence,  $g(z_1) \leq g(x_r)$  for each  $r \geq r_0$  thus one would have

$$\begin{aligned} g(z_2) &< g(z_1) \\ &\leq \inf_r g(x_r) \\ &= \inf_{x \in I} g(x), \end{aligned}$$

which is a contradiction. Hence,  $g(z_2) \geq g(z_1)$ , which shows 1.

To show 2, by the quasiconvexity of  $g$  it is enough to show that, if  $g(z_1) = g(z_2)$ , then  $\{z_1, z_2\} \subset \arg \min_{x \in I} g(x)$ . Suppose that, on the contrary,  $g(z_1) = g(z_2) > \inf_{x \in I} g(x)$ . Then, by Lemma 7,

$$\begin{aligned} g(z_1) &= g(z_2) \\ &> \text{cl}_1 g(x^*), \end{aligned}$$

and we could take a sequence  $\{x_r\}$  converging to  $x^*$  with  $g(x_r)$  converging to  $\text{cl}_1 g(x^*)$  and  $g(x_r) < g(z_2)$  for each  $r$ . Since  $z_1 \in ]x^*, z_2[$ , it would follow that  $z_1 \in ]x_r, z_2[$  for some  $r$ , thus, by the strict quasiconvexity of  $g$ ,  $g(z_1) < g(z_2)$ , which would be a contradiction. Hence  $g(z_1) = g(z_2) = \text{cl}_1 g(x^*)$ , showing that

$$[z_1, z_2] \subset \arg \min_{x \in I} g(x).$$

By the quasiconvexity of both  $g$  and  $\text{cl}_1 g$ , and the optimality of  $x^*$  and  $[z_1, z_2]$  for  $\min_{x \in I} \text{cl}_1 g(x)$ , it then follows that

$$[x^*, z_1] \subset \arg \min_{x \in I} g(x),$$

and the result holds.  $\square$



Another interesting property, which will be exploited in the sequel, states that, once problem  $\min_{x \in I} \text{cl}_I g(x)$  has been solved, any problem  $\inf_{x \in J} g(x)$  with nested feasible interval  $J \subset I$  is immediately solved. Indeed, denoting by  $i(J)$  the interior of  $J$ , one has:

**PROPOSITION 9.** *Let  $J := [a, b] \subset I$  be two compact intervals in  $\mathbb{R}$ . One has:*

1.  $\text{cl}_I g \leq \text{cl}_J g$  on  $J$ , and

$$\text{cl}_I g(x) = \text{cl}_J g(x) \quad \text{for all } x \in i(J) \quad (3.4)$$

2. If  $(\arg \min_{x \in I} \text{cl}_I g(x)) \cap i(J) \neq \emptyset$ , then

$$\inf_{x \in J} g(x) = \min_{x \in I} \text{cl}_I g(x) \quad (3.5)$$

3. If  $(\arg \min_{x \in I} \text{cl}_I g(x)) \cap i(J) = \emptyset$ , then

$$\inf_{x \in J} g(x) = \min\{g(a), g(b)\} \quad (3.6)$$

*Proof.* Part 1 is a direct consequence of the definition of the closure of  $g$  and Lemma 7.

For Part 2, let  $x^* \in \arg \min_{x \in I} \text{cl}_I g(x) \cap i(J)$ ; then, by Parts 1, 2 of Lemma 7 and Part 1 of this proposition,

$$\begin{aligned} \min_{x \in I} \text{cl}_I g(x) &= \text{cl}_I g(x^*) \\ &= \text{cl}_J g(x^*) \\ &= \min_{x \in J} \text{cl}_J g(x) \\ &= \inf_{x \in J} g(x) \\ &\geq \inf_{x \in I} g(x) \\ &= \min_{x \in I} \text{cl}_I g(x) \end{aligned}$$

Part 3 immediately follows from Lemma 8 if  $\arg \min_{x \in I} \text{cl}_I g(x)$  contains points in  $I \setminus J$ . In the remaining case,  $\arg \min_{x \in I} \text{cl}_I g(x)$  consists of just one endpoint of  $J$ , say  $a$ . If a sequence  $\{x_i\} \subset J$  exists converging to  $a$  with  $g(x_i)$  converging to  $\min_{x \in I} \text{cl}_I g(x) = \text{cl}_I g(a)$ , then the result follows from the definition of  $\text{cl}_I g$ . Otherwise there exists  $x^* < a$  with  $g(x^*) < g(a)$  and then the quasiconvexity of  $g$  implies that, for any  $x \in J$ ,

$$\begin{aligned} g(x^*) &< g(a) \\ &\leq \max\{g(a), g(x)\}, \end{aligned}$$

thus  $g(x) \geq g(a)$ , showing (3.6).  $\square$

## 3.2. MULTIPLE-OBJECTIVE QUASICONVEX MAXIMIZATION PROBLEMS ON AN INTERVAL

In this subsection we show how to find a (weak) minimal dominator for the problem  $(P[F; I])$  when  $I = [a, b]$  is a compact interval of  $\mathbb{R}$ .

By Lemma 7, for each  $i = 1, 2, \dots, k$ , the set  $\arg \min_{x \in I} \text{cl}_1 F_i(x)$  is a nonempty closed subinterval of  $I$ , thus it has the form  $[\alpha^i, \beta^i]$ .

**LEMMA 10.** *Let  $x \in ]a, \min_{1 \leq i \leq k} \beta^i[$  (respectively  $x \in ]\max_{1 \leq i \leq k} \alpha^i, b[$ ). Then,  $a \in \mathcal{S}^\geq(x)$  (respectively  $b \in \mathcal{S}^\geq(x)$ ).*

*Proof.* Given  $x \in ]a, \min_{1 \leq i \leq k} \beta^i[$  and  $j \in \{1, 2, \dots, k\}$ , it follows that  $x < \beta^j$ , thus, by the definition of  $\beta^j$  there exists  $y^j \in \arg \min_{y \in I} \text{cl}_1 F_j(y)$  such that  $x \in ]a, y^j[$ . Hence, by Lemma 8,  $F_j(x) \leq F_j(a)$  for all  $j$ , showing that  $a \in \mathcal{S}^\geq(x)$ . The other case is similar.

**PROPOSITION 11.** *Define  $D_I^0$  as*

$$D_I^0 = \{a, b\} \cup \left[ \min_{1 \leq i \leq k} \beta^i, \max_{1 \leq i \leq k} \alpha^i \right],$$

where it is understood that  $[\min_{1 \leq i \leq k} \beta^i, \max_{1 \leq i \leq k} \alpha^i] = \emptyset$  if  $\min_{1 \leq i \leq k} \beta^i > \max_{1 \leq i \leq k} \alpha^i$ . Define also

$$D_I = \begin{cases} \{\alpha\}, & \text{if } F(a) \geq F(b) \\ \{b\}, & \text{if } F(b) \geq F(a), F(b) \neq F(a) \\ D_I^0 \setminus (\{x \neq a : a \in \mathcal{S}^\geq(x)\} \cup \{x \neq b : b \in \mathcal{S}^\geq(x)\}), & \text{otherwise} \end{cases}$$

One then has

1.  $D_I^0 \in \mathcal{D}[F; I]$ .
2.  $D_I \in \mathcal{D}_{\text{WM}}[F; I]$ .
3. If  $a \in \mathcal{S}^\geq(b)$ ,  $b \in \mathcal{S}^\geq(a)$ , or each  $F_i$  is semistrictly quasiconvex on  $[a, b]$ , then  $D_I \in \mathcal{D}_M[F; I]$ .

*Proof.* Part 1 follows from Lemma 10. To show 2, if  $a \in \mathcal{S}^\geq(b)$  one would have for each  $x \in I$  and  $i \in \{1, 2, \dots, k\}$  that

$$\begin{aligned} F_i(x) &\leq \max\{F_i(a), F_i(b)\} \\ &= F_i(a), \end{aligned}$$

thus  $a \in \mathcal{S}^\geq(x)$ ; hence  $D_I = \{a\} \in \mathcal{D}[F; I]$ , which is (weak) minimal being a singleton. A similar result is obtained when  $b \in \mathcal{S}^\geq(a)$ , thus to finish the proof of 2, we assume that  $a \notin \mathcal{S}^\geq(b)$  and  $b \notin \mathcal{S}^\geq(a)$ . In particular,  $\{a, b\} \subset D_I$ . Given  $x \in [a, b]$ , it follows from Part 1 that there exists some  $y \in \mathcal{S}^\geq(x) \cap D_I^0$ ; if  $y \notin D_I$ , then  $y \notin \{a, b\}$  and either  $a \in \mathcal{S}^\geq(y) \subset \mathcal{S}^\geq(x)$  or  $b \in \mathcal{S}^\geq(y) \subset \mathcal{S}^\geq(x)$ , hence  $\emptyset \neq \{a, b\} \cap \mathcal{S}^\geq(x) \subset D_I \cap \mathcal{S}^\geq(x)$ ; if  $y \in D_I$  then  $D_I \cap \mathcal{S}^\geq(x) \neq \emptyset$ . Thus  $D_I \in \mathcal{D}[F; S]$ .

To show that  $D_I \in \mathcal{D}_{\text{WM}}[F; I]$ , suppose that, by contradiction,  $x, y \in D_I$  exist such

that  $x \in \mathcal{S}^>(y)$ . Since either  $x \in [a, y[$  or  $x \in ]y, b]$ , we can assume w.l.o.g. that  $x \in [a, y[$ . Then, for each  $i$

$$\begin{aligned} F_i(y) &< F_i(x) \\ &\leq \max\{F_i(a), F_i(y)\} \end{aligned}$$

thus  $F_i(y) < F_i(a)$  for each  $i$ . Hence  $a \in \mathcal{S}^>(y)$ , thus  $y \notin D_I$ , which is a contradiction. Hence  $D_I \in \mathcal{D}_{WM}[F; I]$ , and this shows 2.

The minimality property of Part 3 was shown above for  $a \in \mathcal{S}^\geq(b)$  or  $b \in \mathcal{S}^\geq(a)$ , so we show now the case of semistrictly quasiconvex functions  $F_i$ . Suppose that, on the contrary,  $x, y \in D_I$  exist such that  $x \in \mathcal{S}^\geq(y) \setminus \{y\}$ . Since  $y \in D_I$ , one gets that  $x \notin \{a, b\}$ ; then,  $x \in ]a, y[ \cup ]y, b[$ , thus w.l.o.g. we assume  $x \in ]a, y[$ . Since  $x \in D_I \setminus \{a\}$ ,  $a \notin \mathcal{S}(x)$ , thus there exists some  $i$  with  $F_i(a) < F_i(x)$ , thus

$$\begin{aligned} F_i(a) &< F_i(x) \\ &\leq \max\{F_i(a), F_i(y)\}, \end{aligned}$$

thus  $F_i(x) \leq F_i(y)$ , and, since  $x \in \mathcal{S}^\geq(y)$ ,  $F_i(x) = F_i(y)$ , which contradicts the semistrict quasiconvexity of  $F_i$ . Hence,  $D_I \in \mathcal{D}_M[F; I]$ .

**REMARK 12.** In Part 1 of Proposition 11, a dominator has been constructed, consisting of at most three intervals, two of which are reduced to a point. Moreover, such a dominator is easily derived once all the single-objective one-dimensional problems  $\min_{x \in I} \text{cl}_I F_i(x)$ ,  $i = 1, 2, \dots, k$  have been solved.

On the other hand, it follows from the quasiconvexity of the functions  $F_i$  that the set  $\{x \in I : a \in \mathcal{S}^\geq(x)\}$  (respectively  $\{x \in I : b \in \mathcal{S}^\geq(x)\}$ ) is an interval having  $a$  (respectively  $b$ ) as one of its endpoints. This implies that the set  $D_I$ , shown in Part 2 of Proposition 11 to be weak minimal, also consists of at most three intervals, two of which are reduced to the endpoints of  $I$ .

In the case of continuous  $F_i$ , finding the set  $\{x : a \in \mathcal{S}^\geq(x)\}$  is reduced to finding, for each  $i = 1, \dots, k$ , the highest root of the nonlinear equation  $F_i(x) = F_i(a)$ , which, due to the quasiconvexity of  $F_i$  can be solved with any prespecified accuracy by e.g. binary search.

**REMARK 13.** For the biobjective case ( $k = 2$ ), the interval  $[\min_k \beta^k, \max_k \alpha^k]$  is, by construction, such that, within it, both  $F_1$  and  $F_2$  are monotonic: one non-decreasing and the other nonincreasing. Hence, for the biobjective case, there is no loss of generality in assuming that functions  $F_i$  are not only quasiconvex but also quasiconcave on the intervals  $[\min_k \beta^k, \max_k \alpha^k]$ .

**EXAMPLE 1.** Let  $I = [0, 4]$ , and consider the three quasiconvex functions  $F_1, F_2, F_3$  defined as

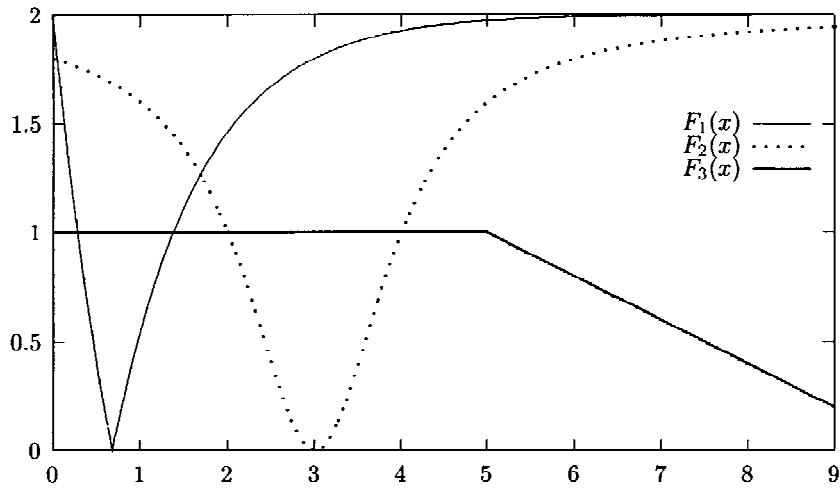


Figure 3. Functions of Example 1.

$$F_1(x) = |4e^{-x} - 2|$$

$$F_2(x) = \frac{2(x-3)^2}{1+(x-3)^2}$$

$$F_3(x) = \min\left\{1, 2 - \frac{x}{5}\right\},$$

depicted in Figure 3.

In order to construct the dominator(s) described in Proposition 11, we must determine first the set  $[\alpha_i, \beta_i]$  of minima on  $I$  for each  $F_i$ . These are respectively  $\{\ln 2\} = \{0.6931\}$ ,  $\{3\}$  and  $[0, 4]$ . This yields

| $i$ | $\alpha_i$ | $\beta_i$ |
|-----|------------|-----------|
| 1   | 0.6931     | 0.6931    |
| 2   | 3          | 3         |
| 3   | 0          | 4         |

For this we obtain the dominator

$$D_I^0 = \{0, 4\} \cup [0.6931, 3]$$

Moreover, by comparing the endpoints, we get

$$F(0) = (2, 1.8, 1)$$

$$F(4) = (1.9267, 1, 1),$$

thus  $F(0) \geq F(4)$ . Hence, by Proposition 11, the set  $D_I = \{0\}$  is not only a weak minimal dominator but also a minimal dominator.

Suppose now that the feasible region is the interval  $I = [5, 9]$ . In this case we obtain

| $i$ | $\alpha_i$ | $\beta_i$ |
|-----|------------|-----------|
| 1   | 5          | 5         |
| 2   | 5          | 5         |
| 3   | 9          | 9         |

From this it is easily seen that

$$D_I^0 = D_I = [5, 9].$$

Since all the functions are semistrictly quasiconvex in  $I$ , it follows that  $I$  is a minimal dominator for  $\max_{x \in I} F(x)$ .

Finally, for  $I = [4, 9]$  we similarly obtain  $D_I^0 = D_I = [4, 9]$ , but in this case  $D_I^0$  is not a minimal dominator, since  $[5, 9]$  is a strictly included dominator (which may be seen to be minimal).

### 3.3. MULTI-OBJECTIVE QUASICONVEX MAXIMIZATION PROBLEMS ON A SET OF INTERVALS

As a natural extension of the model presented in Section 3.2, we address here the problem

$$\max_{x \in X} F(x),$$

where

- $X = \bigcup_{1 \leq i \leq t} I_i$ , with  $\{I_i\}_{1 \leq i \leq t}$  being a family of compact (possibly degenerate) intervals of the real line, not necessarily disjoint,
- $F_1, \dots, F_k$  are quasiconvex on each  $I_i$ ,  $i = 1, \dots, t$ . (Note that this is a weaker assumption than each component of  $F$  to be quasiconvex in the convex hull of  $\bigcup_{1 \leq i \leq t} I_i$ ).

By Proposition 2, if one finds, for each  $i = 1, 2, \dots, t$  some dominator  $D_i \in \mathcal{D}[F; I_i]$ , then any  $D \in \mathcal{D}[F; \bigcup_{1 \leq i \leq t} D_i]$  would serve as a dominator for  $(P[F; \bigcup_{1 \leq i \leq t} I_i])$ . Moreover, if a (weak) minimal dominator is sought, redundant alternatives should be purged, either in the construction of the sets  $D_i$  (by imposing e.g.  $D_i \in \mathcal{D}_{\mathcal{WM}}[F; I_i]$ ) or when they are merged to produce a (small) final dominator.

To approximate this goal one can use a Branch-and-Bound scheme, similar to the one described in [14]: we start with a list  $\mathcal{L}$  of compact intervals, the union of which is known to be a dominator for  $(P[F; \bigcup_{1 \leq i \leq t} I_i])$ , and then refine iteratively

the elements in  $\mathcal{L}$ , by making pairwise comparisons, in such a way that, at any stage, one has

$$\bigcup_{I \in \mathcal{L}} I \in \mathcal{D}[F; \bigcup_{1 \leq i \leq l} I_i]$$

To perform comparisons among elements in  $\mathcal{L}$  we introduce, for each interval  $I := [a, b]$  contained in some  $I_i$ , the vectors  $M(I), UB(I) \in \mathbb{R}^k$  of evaluations at the midpoint of  $I$  and a componentwise upper bound of  $F$ , respectively:

$$M(I)_j = F_j\left(\frac{a+b}{2}\right)$$

$$UB(I)_j \geq \max_{x \in I} F_j(x)$$

REMARK 14. By the quasiconvexity of  $F_j$  on  $I \subset I_i$ , it follows that one may choose

$$UB(I)_j = \max\{F_j(a), F_j(b)\} \quad j = 1, 2, \dots, k$$

Note also that for  $I = \{a\}$  we have  $M(I) = UB(I)$ .

From the definitions of the vectors  $M$  and  $UB$  one immediately obtains the following way to check whether some interval  $J$  can be discarded from further consideration in the Branch-and-Bound scheme.

PROPOSITION 15. *Given nonempty compact intervals  $I, J$ , suppose  $F$  is continuous on  $I$  and on  $J$ . Then the following statements are equivalent:*

1.  $I \in \mathcal{D}[F; J]$ , i.e. for any  $y \in J$  there exists  $x \in I$  with  $F(x) \geq F(y)$
2.  $0 \leq \min_{y \in J} \max_{x \in I} \min_{1 \leq j \leq k} (F_j(x) - F_j(y))$ .

This is implied by both

$$\bigcap_{1 \leq j \leq k} \{x \in I : F_j(x) \geq UB(J)_j\} \neq \emptyset \quad (3.7)$$

and

$$0 \leq \min_{1 \leq j \leq k} (M(I)_j - UB(J)_j), \quad (3.8)$$

while (3.8) always implies (3.7).

*Proof.* The equivalence between 1 and 2 is evident. Since (3.7) is equivalent to the existence of some  $x \in I$  with

$$F(x) \geq F(y) \quad \forall y \in J, \quad (3.9)$$

it clearly implies 1. On the other hand, (3.8) is equivalent to (3.9) for  $x$  fixed to the midpoint of  $I$ . Hence, (3.8) implies (3.7) and the result follows.  $\square$

Although condition (3.8) is easier to implement, the stronger test (3.7) is also of practical interest since this intersection set, if nonempty, has a simple structure due to the quasiconvexity of  $F$ , as indicated by the following simple result:

PROPOSITION 16. *One has for any values  $c_j$*

1. *Each set  $\{x \in I : F_j(x) \geq c_j\}$  consists of at most two intervals, each with an endpoint of  $I$  as one of its endpoints.*
2. *For  $k_0 = 1, 2, \dots, k$ , the set  $\bigcap_{1 \leq j \leq k_0} \{x \in I : F_j(x) \geq c_j\}$  is a collection of  $n(k_0)$  intervals, with*

$$n(1) \leq 2$$

$$n(k_0) \leq n(k_0 - 1) + 1, \quad k_0 = 2, 3, \dots, k$$

The basic steps of the Branch-and-Bound procedure are described below:

**Algorithm 1**

**Initialization:**  
 Set  $\mathcal{L} := \{cl(D_{I_j}), j = 1, \dots, t\}$   
 Set  $r := 1$

**Iteration  $r = 1, 2, \dots, :$**   
 for all  $I \in \mathcal{L}$  do  
   If, for some  $J \in \mathcal{L}, J \neq I$ , (3.8) or (3.7) hold, then  
   delete  $I$  from  $\mathcal{L}$ ;  
   Else, if  $I$  is non-degenerate do  
   split  $I$  into  $I_1$  and  $I_2$  at the midpoint of  $I$ ;  
   replace  $I$  by  $I_1$  and  $I_2$  in  $\mathcal{L}$ ;  
 GoTo Iteration  $r + 1$

Before discussing the output of the algorithm in the limit ( $r = \infty$ ), let us present an illustrative example.

EXAMPLE 1 (Cont.)

Let  $F$  be the three-objective one-dimensional function described in the first part of the Example, and assume now that the feasible region  $X$  consists of the two compact segments  $I_1 = [0, 4]$ , and  $I_2 = [5, 9]$ .

In the Initialization phase, we must construct the sets  $cl(D_{I_j}), j = 1, 2$ . This was already done in the first part of the Example, thus we start with the list

$$\mathcal{L} = \{\{0\}, [5, 9]\}.$$

Then, we go to Iteration 1. For each interval  $I$  (degenerate or not) in  $\mathcal{L}$ , the vectors  $M(I), UB(I)$  must be constructed. (Observe that this task becomes trivial using Remark 14 above.) Evaluations at the endpoints, 0, 5, 9 and the midpoint 7 yield

$$F(0) = (2, 1.8000, 1)$$

$$F(5) = (1.9730, 1.6000, 1)$$

$$F(7) = (1.9964, 1.8824, 0.6000)$$

$$F(9) = (1.9995, 1.9459, 0.2000)$$

We then obtain

| $I$      | $M(I)$                     | $UB(I)$               |
|----------|----------------------------|-----------------------|
| $\{0\}$  | $(2, 1.8000, 1)$           | $(2, 1.8000, 1)$      |
| $[5, 9]$ | $(1.9964, 1.8824, 0.6000)$ | $(1.9995, 1.9459, 1)$ |

We will only use the simplest test, namely, (3.8) in the algorithm.

Since no pair of intervals in  $\mathcal{L}$  satisfies condition (3.8), we go to Iteration 2 with the list of intervals

$$\mathcal{L} = \{\{0\}, [5, 7], [7, 9]\}.$$

Two new midpoints appear, namely, 6 and 8, with objective values

$$F(6) = (1.9901, 1.8000, 0.8000)$$

$$F(8) = (1.9987, 1.9231, 0.4000).$$

This enables us to update the table of vectors  $M, UB$  yielding

| $I$      | $M(I)$                     | $UB(I)$                    |
|----------|----------------------------|----------------------------|
| $\{0\}$  | $(2, 1.8000, 1)$           | $(2, 1.8000, 1)$           |
| $[5, 7]$ | $(1.9901, 1.8000, 0.8000)$ | $(1.9964, 1.8824, 1)$      |
| $[7, 9]$ | $(1.9987, 1.9231, 0.4000)$ | $(1.9995, 1.9459, 0.6000)$ |

As in the previous iteration, no pair of intervals satisfies condition (3.8), and we go to Iteration 3 with the updated list of intervals

$$\mathcal{L} = \{\{0\}, [5, 6], [6, 7], [7, 8], [8, 9]\}$$

The new midpoints give objective values

$$F(5.5) = (1.9837, 1.7241, 0.9000)$$

$$F(6.5) = (1.9940, 1.8491, 0.7000)$$

$$F(7.5) = (1.9978, 1.9059, 0.5000)$$

$$F(8.5) = (1.9992, 1.9360, 0.3000)$$



With this, our new table of vectors  $M, UB$  is given by

| $I$     | $M(I)$                   | $UB(I)$                  |
|---------|--------------------------|--------------------------|
| $\{0\}$ | (2, 1.8000, 1)           | (2, 1.8000, 1)           |
| [5, 6]  | (1.9837, 1.7241, 0.9000) | (1.9901, 1.8000, 1)      |
| [6, 7]  | (1.9940, 1.8491, 0.7000) | (1.9964, 1.8824, 0.8000) |
| [7, 8]  | (1.9978, 1.9059, 0.5000) | (1.9987, 1.9231, 0.6000) |
| [8, 9]  | (1.9992, 1.9360, 0.3000) | (1.9995, 1.9459, 0.4000) |

In this case, the sufficient condition for dominance is satisfied for the pair of intervals  $\{0\}$  and  $[5, 6]$ , so the interval  $[5, 6]$  can be excluded for further considerations.

We would then obtain a reduced list

$$\mathcal{L} = \{\{0\}, [6, 7], [7, 8], [8, 9]\}$$

to start Iteration 4, if desired. □

The following theorem shows that the successive steps of the algorithm above provide a sequence of nested compact dominators, converging to a dominator which, under mild further assumptions on the functions  $F_i$ , enjoys minimality properties:

**PROPOSITION 17.** *Denote by  $D_r$  the union of all intervals of  $\mathcal{L}$  at the end of iteration  $r$ , and by  $D^*$  the compact set*

$$D^* = \bigcap_{r=1}^{\infty} D_r$$

1.  $D_1 = X = \bigcup_{1 \leq i \leq t} I_i$  and  $D_{r+1} \subset D_r$  for all  $r$ .
2. If  $F$  is upper-semicontinuous, then

$$D^* \in \mathcal{D}[F; X] \tag{3.10}$$

3. Moreover, if  $F$  is continuous, then

$$D^* \in \mathcal{D}_{w.u.}[F; X]. \tag{3.11}$$

*Proof.* The first property is evident from the algorithm.

By construction, each  $D_r$  is compact, thus their intersection is also compact. Moreover,  $D_r \in \mathcal{D}[F; X]$ , thus, by Proposition 3, (3.10) follows.

To show (3.11), suppose, on the contrary, that there exist  $x_1, x_2 \in D^*$  with  $x_1 \in \mathcal{S}^>(x_2)$ . If, for each  $i = 1, 2$  and  $r = 1, 2, \dots$ , we denote by  $\mathcal{I}_i^r$  the class of intervals  $I_i^r$  in the list at stage  $r$  with  $x_i \in I_i^r$ , it will follow from the splitting process that there exists some  $r_0$  such that, for each  $r \geq r_0$ , and each  $I_i^r \in \mathcal{I}_i^r$

$$x_1 \notin I_2^r, \quad \text{and} \quad x_2 \notin I_1^r$$

Since the functions  $F_i$  are continuous, thus uniformly continuous on  $X$ , there would exist some  $r$  such that for each  $I_i^r \in \mathcal{I}_i^r$

$$F_j(x) > F_j(y) \quad \text{for all } x \in I_1^r \text{ and } y \in I_2^r, \quad j = 1, 2, \dots, k$$

Hence  $UB(I_2^r) < M(I_1^r)$ , implying that  $I_2^r$  (thus  $x_2$ ) would have been deleted prior to stage  $r$  by (3.8), thus  $x_2 \notin D^*$ , which is a contradiction.  $\square$

#### 4. Multiple-objective multi-dimensional problems

For the single-objective case (i.e., if  $k = 1$  in  $(P[F_1; S])$ ), it is a well-known result of Global Optimization that, if  $S$  is a polytope and  $F_1$  is quasiconvex on  $S$ , then the set of vertices of  $S$  is a dominator for  $(P[F_1; S])$ , [15].

In other words, if, for  $j = 0, 1, \dots, n$ ,  $\mathcal{F}^j$  denotes the set of points of a polytope  $S$  contained in some  $j$ -dimensional face of  $S$ , then

$$\mathcal{F}^0 \in \mathcal{D}[F_1; S] \tag{4.12}$$

The next proposition extends assertion (4.12) to multiple-objective quasiconvex problems. To show it, we will use the following

LEMMA 18. *Let  $P$  be a polyhedron in  $\mathbb{R}^n$ , and let  $H_1, H_2, \dots, H_t$  be closed halfspaces in  $\mathbb{R}^n$ . If  $x^*$  is an extreme point of  $P \cap \bigcap_{1 \leq i \leq t} H_i$ , then  $x^*$  belongs to some face of  $P$  with dimension not greater than  $t$ .*

*Proof.* Let  $P$  be represented as

$$P = \{x \in \mathbb{R}^n : a'_r x \leq b_r \text{ for all } r \in R\}$$

for some finite index set  $R$ , and let each  $H_i$  be given as

$$\{x \in \mathbb{R}^n : c'_i x \leq d_i\}$$

Define the sets of active indices  $R(x^*)$  and  $T(x^*)$  as

$$R(x^*) = \{r \in R : a'_r x^* = b_r\}$$

$$T(x^*) = \{i, 1 \leq i \leq t : c'_i x^* = d_i\}$$

Then  $x^*$  belongs to the face  $F$  of  $P$ ,

$$F = P \cap \{x \in \mathbb{R}^n : a'_r x = b_r, \forall r \in R(x^*)\}$$

We will show that  $F$  has dimension not greater than  $t$ . Indeed, since  $x^*$  is, by assumption, an extreme point of  $P \cap \bigcap_{1 \leq i \leq t} H_i$ , then the set of vectors  $\{a_r\}_{r \in R(x^*)} \cup \{c_i\}_{i \in T(x^*)}$  has rank

$$\text{rank}(\{a_r\}_{r \in R(x^*)} \cup \{c_i\}_{i \in T(x^*)}) = n$$

Hence, denoting by  $|T(x^*)|$  the cardinality of  $T(x^*)$ , one obtains

$$\begin{aligned} \text{rank}(\{a_r\}_{r \in R(x^*)}) &\geq n - |T(x^*)| \\ &\geq n - t, \end{aligned}$$

thus the dimension of  $F$  cannot be greater than  $t$ .  $\square$

**PROPOSITION 19.** *Let  $S$  be a polytope in  $\mathbb{R}^n$ , let  $k \leq n + 1$ , and let  $F_1, \dots, F_k$  be quasiconvex functions on  $S$ , all but possibly one of which are upper-semicontinuous. Then*

$$\mathcal{F}^{k-1} \in \mathcal{D}[F; S] \quad (4.13)$$

*Proof.* Without loss of generality we assume that  $F_1, F_2, \dots, F_{k-1}$  are upper-semicontinuous on  $S$ . We will show that, for any  $x \in S$ ,

$$\mathcal{S}^\geq(x) \cap \mathcal{F}^{k-1} \neq \emptyset \quad (4.14)$$

Let  $x \in S$ , and denote by  $\mathcal{A}(x)$  the index set

$$\mathcal{A}(x) = \{i, 1 \leq i \leq k-1, F_i(y) < F_i(x) \text{ for some } y \in S\}.$$

If  $\mathcal{A}(x)$  is empty, we would have

$$F(y) \geq F(x) \quad \forall y \in S,$$

thus any vertex  $y^*$  of  $S$  satisfies  $y^* \in \mathcal{S}^\geq(x)$ . Hence

$$\emptyset \neq \mathcal{S}^\geq(x) \cap \mathcal{F}^0 \subset \mathcal{S}^\geq(x) \cap \mathcal{F}^{k-1},$$

showing (4.14).

We consider now the case  $\mathcal{A}(x) \neq \emptyset$ . For each  $i \in \mathcal{A}(x)$ , the convex set  $\{y \in S : F_i(y) < F_i(x)\}$  is open in  $S$  (its complement is closed due to the upper-semicontinuity of  $F_i$ ) and does not contain  $x$ . Hence, there exists some nonzero vector  $u^i$  such that

$$\langle u^i, y - x \rangle > 0 \quad \text{for all } y \in S \text{ with } F_i(y) < F_i(x), \quad (4.15)$$

where  $\langle \cdot, \cdot \rangle$  stands for the usual scalar product in  $\mathbb{R}^n$ .

Consider the polyhedron  $S(x)$ ,

$$S(x) = S \cap \{y \in \mathbb{R}^n : \langle u^i, y - x \rangle \leq 0, \quad \forall i \in \mathcal{A}(x)\},$$

which is nonempty because  $x \in S(x)$ . Consider the optimization problem

$$\max_{y \in S(x)} F_k(y) \quad (4.16)$$

Since  $F_k$  is quasiconvex on the nonempty polyhedron  $S(x)$ , (4.16) has an optimal solution at some vertex  $y^*$  of  $S(x)$ . We will show that

$$y^* \in \mathcal{F}^{k-1} \cap \mathcal{S}^\geq(x) \quad (4.17)$$

Since  $y^*$  is a vertex of  $S(x)$ , Lemma 18 implies that

$$y^* \in \mathcal{F}^{|\mathcal{A}(x)|} \subset \mathcal{F}^{k-1} \quad (4.18)$$

Since  $x \in S(x)$  and  $y^*$  is optimal for (4.16),

$$F_k(y^*) \geq F_k(x) \quad (4.19)$$

By definition of  $\mathcal{A}(x)$ ,

$$F_i(y^*) \geq F_i(x) \quad \forall i \in \{1, 2, \dots, k-1\} \setminus \mathcal{A}(x) \quad (4.20)$$

and by (4.15) and the fact that  $y^* \in S(x)$ ,

$$F_i(y^*) \geq F_i(x) \quad \forall i \in \mathcal{A}(x) \quad (4.21)$$

Joining (4.18–4.21), (4.17) holds, thus  $\mathcal{S}^{\geq}(x) \cap \mathcal{F}^{k-1} \neq \emptyset$ , as asserted.  $\square$

REMARK 20. The assumption of upper-semicontinuity of at least  $k-1$  functions is not superfluous, as the following example shows: let  $k=2$ ,  $n=2$ ,  $S=[0,1] \times [0,1]$ , and the functions  $F_1, F_2$  defined as

$$F_1(x) = \begin{cases} 0, & \text{if } x_2 > \frac{1}{2} \text{ or } x = (0, \frac{1}{2}) \\ 1, & \text{otherwise} \end{cases} \quad F_2(x) = \begin{cases} 0, & \text{if } x_2 > \frac{1}{2} \text{ or } x = (1, \frac{1}{2}) \\ 1, & \text{otherwise} \end{cases}$$

Both functions are quasiconvex but are not upper-semicontinuous; let  $x^* = (\frac{1}{2}, \frac{1}{2})$ . It is easily seen that

$$\mathcal{S}^{\geq}(x^*) = \{(\lambda, \frac{1}{2}) : 0 < \lambda < 1\}$$

thus  $\mathcal{S}^{\geq}(x^*) \cap \mathcal{F}^1 = \emptyset$ , showing that  $\mathcal{F}^1$  is not a dominator.  $\square$

As a consequence of Propositions 19 and 2, one obtains

COROLLARY 21. *Let  $S$  be the union of  $t$  polytopes  $S_1, \dots, S_t$  in  $\mathbb{R}^n$ . Let  $F_1, \dots, F_k$  be  $k \leq n+1$  real-valued functions on  $S$ . On each  $S_j$ , let all  $F_i$  be quasiconvex and all but possibly one  $F_i$  be lower-semicontinuous. Then the union of all  $k-1$ -faces of all  $S_j$  is a dominator for  $P[F; S]$ .*

Proposition 19 also enables us to derive localization results for single-objective problems.

COROLLARY 22. *Let  $S$  be the union of  $t$  polytopes  $S_1, \dots, S_t$  in  $\mathbb{R}^n$ . Let  $F_1, \dots, F_k$  be  $k \leq n+1$  real-valued functions on  $S$ , quasiconvex on each  $S_j$ . For any componentwise nondecreasing  $\Phi : F(S) \rightarrow \mathbb{R}$  such that Problem*

$$\max_{x \in S} \Phi(F_1(x), F_2(x), \dots, F_k(x))$$

has an optimal solution, the union of the set of  $k - 1$ -faces of all  $S_j$  also contains an optimal solution.

In particular, for  $F_1, F_2, \dots, F_k$  linear fractional functions with positive denominators on a polytope  $S$ , which are well-known to be quasiconvex (see e.g. [1], p. 165) and  $\Phi(s_1, s_2, \dots, s_k) = s_1 + s_2 + \dots + s_k$ , or  $\Phi(s_1, s_2, \dots, s_k) = \max\{s_1, s_2, \dots, s_k\}$ , we obtain

**COROLLARY 23.** *The minimum of the sum (respect. the maximum) of  $k$  linear fractional functions with positive denominators over a polytope  $S$  in dimension  $n \geq k - 1$  is attained at some  $k - 1$ -face of  $S$ .*

This generalizes the results known for the case  $k = 2$  (see the review of [27] and the references therein). For an application see [25].

**REMARK 24.** For biobjective problems ( $k = 2$ ), since both  $F_j$  are quasiconvex on each edge, after embedding such edges as compact intervals of the real line, one can use the results in Section 3 to design an algorithm converging to a weak minimal dominator.

For the case of general  $k$ , Proposition 19 seems at the moment to be mainly of theoretical interest: In principle, Algorithm 1 can be generalized to the  $k$ -dimensional case, by replacing intervals by e.g. simplices, although the corresponding bounding scheme does not extend to the general case, and less efficient schemes, such as those proposed in [4, 14], should be used.

Nevertheless, this kind of localization results can be used to design new heuristic resolution methods of problems of the form  $\min_{x \in S} \Phi(F(x))$ , where  $k$ , the number of components of  $F$ , is very small, and, in particular, much smaller than the dimension  $n$  of the space.

We know then that the search for optimal solutions can be reduced to the  $k - 1$ -dimensional faces of  $S$ , so that algorithms which alternate a global search in a given low-dimensional face with moves to adjacent low-dimensional faces, can be used.

## 5. Application: Location of a semi-obnoxious facility

Let  $S = S_1 \cup S_2 \cup \dots \cup S_t$ , each  $S_i$  being a convex polygon in  $\mathbb{R}^2$ . Two finite subsets  $\mathcal{A}^+, \mathcal{A}^-$  of  $\mathbb{R}^2$  are given. Associated with each  $a \in \mathcal{A}^+$  we have a concave function  $g_a : [0, +\infty) \rightarrow \mathbb{R}$  and a polyhedral gauge  $\gamma_a$ , [9, 10, 19], i.e., a Minkowski functional whose unit ball is a polytope.

Let  $h : [0, +\infty) \rightarrow \mathbb{R}$  be a nonincreasing function, and consider the biobjective problem

$$\min_{x \in S} (F_1(x), F_2(x)), \tag{5.22}$$

where

$$F_1(x) = \sum_{a \in \mathcal{A}^+} g_a(\gamma_a(x - a))$$

$$F_2(x) = \max_{a \in \mathcal{A}^-} h(\|x - a\|),$$

$\|\cdot\|$  being the euclidean norm.

This problem has its motivation in Continuous Location of semidesirable facilities, see [17, 23] for an introduction to Continuous Location in general and [6, 24] for semidesirable facility location models: A facility is to be located within region  $S$ , and will interact with individuals who want the facility close (those in  $\mathcal{A}^+$ ) and others who want the facility far (those in  $\mathcal{A}^-$ ). Interactions with  $\mathcal{A}^+$  provide the first objective in (5.22): the minimization of the total transportation cost  $F_1(x)$ , where transportation cost from  $a \in \mathcal{A}^+$  to  $x$  is given by a concave function  $g_a$  of the distance from  $a$  to  $x$ , the latter measured by the polyhedral gauge  $\gamma_a$ , [29].

On the other hand, interactions of the facility with  $\mathcal{A}^-$  provide the second objective  $F_2$ , which measures the highest damage suffered by points in  $\mathcal{A}^-$ , where the damage suffered by  $a \in \mathcal{A}^-$  is assumed to be given by a nonincreasing function  $h$  of the Euclidean distance from  $a$  to  $x$ , see [11, 24].

In practice, the two objectives of (5.22) are aggregated into a single criterion, yielding a problem of the form

$$\max \Phi \left( - \sum_{a \in \mathcal{A}^+} g_a(\gamma_a(x - a)), \max_{a \in \mathcal{A}^-} h(\|x - a\|) \right), \quad (5.23)$$

[6, 24] for some globalizing  $\Phi$ , and the resulting problem (multimodal, as a rule), can be tackled e.g. by the 2-dimensional Branch and Bound method described in [13]. However, as shown below (Proposition 25), the search of an optimal solution for (5.23) can be restricted to a series of segments, thus (5.23) can be solved by simply using single-variable Global-Optimization techniques, [2, 12], which are usually much faster than their two-variable counterparts.

In order to obtain a dominator for (5.22) one should observe first that, since  $h$  is assumed to be nonincreasing, it suffices to obtain a dominator for problem

$$\max_{x \in S} (-F_1(x), \min_{a \in \mathcal{A}^-} \|x - a\|) \quad (5.24)$$

(in fact, if  $h$  is decreasing, both problems are equivalent). Let us rewrite now (5.24) within our framework. For polyhedral gauges, using the concept of *elementary convex set* of [10], one can obtain a subdivision  $\mathcal{C}$  of the plane into polyhedra in such a way that, within each  $C \in \mathcal{C}$ , each gauge  $\gamma_a$  is affine, see [9, 10] for further details. For instance, if each  $\gamma_a$  is the  $l_1$  norm, then the polyhedral subdivision of the plane is obtained after constructing horizontal and vertical lines through each  $a \in \mathcal{A}^+$ , yielding a total of  $O(|\mathcal{A}^+|^2)$  cells.

Moreover, defining, for each  $a \in \mathcal{A}^-$ , the *Voronoi cell*  $V(a)$  associated with  $a$  as

$$V(a) = \{x \in \mathbb{R}^2 : \|x - a\| \leq \|x - b\| \text{ for all } b \in \mathcal{A}^-\},$$

the class  $\mathcal{V} = \{V(a) : a \in \mathcal{A}^-\}$  also constitutes a polyhedral subdivision of  $\mathbb{R}^2$  in  $O(|\mathcal{A}^-|)$  polyhedra, which can be efficiently constructed in  $O(|\mathcal{A}^-| \log |\mathcal{A}^-|)$ , see e.g. [20, 26].

Consider now the class  $\mathcal{Z}$  of all  $Z$  of the form

$$S_i \cap C \cap V(a)$$

for some  $i, 1 \leq i \leq t$ ,  $C \in \mathcal{C}$  and  $a \in \mathcal{A}^-$  which are nonempty. On each  $Z \in \mathcal{Z}$ , we have that  $-F_1$  is convex (it is the composition of the convex function  $-\sum_{a \in \mathcal{A}^+} g_a$  with the affine functions (within  $Z$ !)  $\gamma_a$ , and  $F_2$  is also convex (recall that, for  $Z \in \mathcal{Z}$  fixed, there exists some  $a^* \in \mathcal{A}^-$  such that  $\min_{a \in \mathcal{A}^-} \|x - a\| = \|x - a^*\|$ ). Hence, rewriting (5.24) as

$$\max_{x \in \bigcup_{Z \in \mathcal{Z}} Z} (-F_1(x), \min_{a \in \mathcal{A}^-} \|x - a\|),$$

we can use Corollary 21 to obtain

**PROPOSITION 25.** *The edges of the sets in  $\mathcal{Z}$  constitute a dominator for Problem (5.22).*

After embedding the edges of polytopes in  $\mathcal{Z}$  as compact intervals of the real line, one can use the algorithm described in Section 3.3 to reduce the size of such dominator, converging (in case of decreasing  $h$ ) to a weak minimal dominator.

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